Introduction to libpari programming

Introduction to libpari programming A tutorial

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libpari C headers

PARI code can be compiled in three ways:

- 1. as a standalone program
- 2. as a loadable module
- 3. directly inside libpari

In the first two cases the headers are included as follow

```
#include <pari/pari.h>
```

in the third case

```
#include "paridecl.h"
```

after all extra system headers have been included. In the first case, PARI needs to be initialized with pari_init before being used.

libpari C types

The PARI library API mostly relies on three C types: long, ulong (short for unsigned long) and GEN.

PARI denotes the number of bits in <code>ulong</code> by <code>BITS_IN_LONG</code>. A GEN \times is a pointer to a data structure representing a PARI object.

x[0] contains the type and the length of the object, which are accessed using typ and lg. The other components can be either codeword or pointers to other GEN (which can contains pointers to other GEN etc.) GEN can have several components that point to the same sub-GEN, but cycles are not allowed.

The GEN types

typ returns one of the following enum values.

Leaf types (all components are codeword)

- t_INT arbitrary precision integers
- t_REAL arbitrary precision real numbers
- t_VECSMALL vectors of long
- t_STR character string
- t_INFINITY $\pm\infty$

Recursive types (some components are pointers to other GENS)

- t_intmod $\mathbb{Z}/n\mathbb{Z}$
- t_FRAC rational numbers
- t_FFELT finite field elt.
- t_COMPLEX complex numbers
- t_PADIC *p*-adic numbers
- t_QUAD quadratic numbers (deprecated)
- t_POLMOD K[X]/T

The GEN types

- t_POL polynomials
- t_SER power series
- t_RFRAC rational function
- t_QFB binary quadratic form
- t_VEC row vector
- t_COL column vector
- t_MAT matrix
- t_list list
- t_CLOSURE **GP functions**
- t_ERROR error context

It is customary to call a GEN of type t_INT a t_INT, etc.

Warning about use of long and ulong

- According to the C standard, ulong are wrapping, that is all operations are done modulo 2^{BITS}_IN_LONG</sup>, but this is not the case for long, where overflow are not defined.
- % and / in C follow FORTRAN semantic and not PARI semantic when the operands are negative: -1%3 = -1.
 PARI provides smodss and umodsu to avoid such problem.
- Immediate constants sometime need to be postfixed with L or UL to avoid confusion with int (especially in variadic functions like mkvecsmalln).
- C int must generally be avoided.

GEN

- typ(x): return the type of x.
- ► lg(x): return the length of x.
- settyp(x,t): set the type of x to t.
- setlg(x,l): set the length of x to l.
- cgetg(I,t); allocate a GEN of length I and type t. on the PARI stack.

t_INT object

t_INT are arbitrary precision relative integers.

- signe(x) : sign of x, 0 is x == 0
- lgefint (x) : actual size in word (can be smaller than lg(x)).
- expi(x) : exponent (logint(x, 2)).

Access to the mantissa words of a t_INT is done using the macro int_W, see the documentation. The sign can be chnaged with setsigne.

Small integers are available as universal objects.

t_INT object

In the API, the operand types are encoded by the letter

- s:long (for "small integer")
- u:ulong
- i:t_INT

For example, for conversion:

- stoi: convert a long to a t_INT
- utoi: convert a ulong to a t_INT
- itos: convert a t_INT to a long
- itou: convert a t_INT to a ulong

Comparing:

- equality: equalii, equaliu, equalis
- equality to 1 or -1: equali1, equalim1
- comparison: cmpii, cmpis, cmpiu cmpsi, cmpui, cmpss, cmpuu: return the sign of x - y as a int.

Operations on t_INT

- addii, addis, addiu, addss, adduu: return the sum (return a t_INT).
- idem with add replaced by sub, mul, mod.
- ▶ negi(x) returns -x, absi return |x|.
- sqri, sqrs, sqru return the square.
- shifti(x,n) shift x of n bits (n can be positive or negative).
- truedvmdii, truedivii, modii euclidean division.
- smodis, smodss: return the remainder as a long.
- umodiu, umodsu: return the remainder as a ulong.
- ▶ gc_INT faster version of gc_GEN for t_INT.
- gc_stoi faster version of gc_GEN(av, stoi(...))
- gc_utoi faster version of gc_GEN(av,utoi(...))

t_REAL

 ${\tt t_REAL}$ are arbitrary precision floating points real numbers

- signe(x) : sign of x, 0 is x == 0
- realprec(x) : precision in bit, always a multiple of BITS_IN_LONG.
- expo(x) : exponent of x

mantissa_real(x, &e) return the mantissa as a t_INT.

The sign can be changed with setsigne, the exponent with setexpo.

The code letter for t_REAL is r. Functions that need to convert integers to t_REALs need an extra argument called prec which is the precision wanted.

- stor(x, prec): convert a long to a t_REAL
- utor(x, prec): convert a ulong to a t_REAL
- itor(x, prec): convert a t_INT to a t_REAL
- rtor(x, prec): convert a t_REAL to a t_REAL with a different precision.

Operations on t_REAL

- equality: equalrr, equalri, equalrs
- comparison: cmprr, cmpri, cmprs, cmpir, cmpsr.
- addrr, addri, addrs, addir, addsr: return the sum (return a t_REAL).
- idem with add replaced by sub, mul, div
- negr(x) returns -x, absr(x) return |x|, sqrr(x) returns x². shiftr(x,n) multiply x by 2ⁿ (n can be positive or negative).
- divri, truedivii, modii
- truncr, floorr, ceilr roundr.

Vectors

Vectors are available in two variant <code>t_VEC</code> and <code>t_COL</code>. Since PARI uses French linear algebra convention, <code>t_COL</code> is often more natural.

To test if a type t is either t_VEC and t_COL, use

is_vec_t(t). if v is a vector, and l=lg(v), then v has I - 1 components, gel(v, 1), ..., gel(v, l-1).

To allocate a vector with n undefined components, do

```
v = cgetg(n+1, t_VEC); or v = cgetg(n+1, t_VEC);
```

t_COL);.

Note than this is not a valid object until all components have been set (by using $gel(v, i) = \dots$).

Vector example

```
GEN fun(long n)
{
    long i;
    GEN v = cgetg(n+1, t_COL);
    for (i = 1; i <= n; i++)
        gel(v,i) = sqru(i);
    return v;
}</pre>
```

Vectors

zerovec(n) and zerocol(n) create a vector of gen_0 that can be filled later. const_vec(n,x) and const_col(n,x) create vectors of x.

Fixed-length short vectors can be created with mkvec(x1), mkvec2(x1,x2), mkvec3(x1,x2,x3), mkvec4(x1,x2,x3,x4), mkvec5(x1,x2,x3,x4,x5), mkvecn(n,x1,...,xn), mkcol(x1), mkcol2(x1,x2), mkcol3(x1,x2,x3), mkcol4(x1,x2,x3,x4), mkcol5(x1,x2,x3,x4,x5). mkcoln(n,x1,...,xn). For example [0,1,2] can be created with mkvec3(gen_0,gen_1,gen_2).

t_MAT

t MAT are represented as vector of t COL of identical length. if m is a t_MAT, and I = Ig(m), then m has I - 1 columns, gel(m, 1), ..., gel(m, 1-1), which have all the same length. Thus the number of row of a matrix with zero columns is not defined. The coefficients of *m* can be accessed with gcoeff(m, i, j) which is a short-hand for qel(qel(m,j),i). To allocate a t_MAT with n undefined colums, do m = cgetg(n+1, t MAT) then set the columns with $qel(v,i) = \ldots$

zeromatcopy(n,m) create a matrix of gen_0 that can be filled later.

Matrix example

```
GEN fun(long n, long m)
  long i, j;
  GEN v = cgetg(m+1, t_MAT);
  for (i = 1; i <= m; i++)
  {
    GEN c = cqetq(n+1, t COL);
    for (j = 1; j \le n; j++)
      qel(c,j) = mulss(i,j);
    gel(v, i) = c;
  }
  return m;
```

t_VECSMALL

t_VECSMALL is a low-level type used for vector of long or ulong depending on the context. If v is a t_VECSMALL and l = lg(v), the components are v[1], ..., v[1-1] in the long case and uel(v, 1), ..., uel(v, 1-1). To allocate a t_VECSMALL with n undefined components, do v = cgetg(n+1, t_VECSMALL); and then set v[1], ..., v[n] or uel(v, 1), ..., uel(v, n).

t_VECSMALL example

```
GEN fun(long n)
{
    long i;
    GEN v = cgetg(n+1, t_VECSMALL);
    for (i = 1; i <= n; i++)
        uel(v,i) = i;
    return v;
}</pre>
```

t_VECSMALL

```
zero_zv(n) creates a vector of 0 that can be filled later.
const_vecsmall(n,x) create vectors of x.
Fixed-length short vectors can be created with
mkvecsmall(x1), mkvecsmall2(x1,x2),
mkvecsmall3(x1,x2,x3), mkvecsmall4(x1,x2,x3,x4),
mkvecsmall5(x1,x2,x3,x4,x5),
mkvecsmalln(n,x1,...,xn).
```

t_POL

 ${\tt t_POL}$ are polynomials.

- signe(x): 0 if x = 0, 1 otherwise.
- varn(x): variable number of x.
- degpol(x): degree of x (-1 if x = 0), degpol(x) = lg(x) - 3.
- lgpol(x):1+degpol(x),lg(x)-2.
- leading_coeff(x): leading coefficient.
- constant_coeff(x): constant coefficient.
- pol_0(v), pol_1(v), pol_x(v): polynomials 0, 1, x in variable v.

The leading coefficient must not be an exact zero. However a polynomial can have signe 0 even if its degree is not -1, if all its coefficients are inexact zero.

If *P* is a t_POL of degree *d*, the coefficients of degree

 $0 \le i \le d$ can be accessed with gel(P,i+2).

The variable number can be set with setvarn. All variables that appears in components of polynomial must have strictly lower priorities than varn(x)

Priority are compared using varncmp(v,w).

t_POL

Creating a t_POL of degree *d* and variable number *v* requires four steps:

```
allocation P = cgetg(d+3, t_POL);
settting the variable P[1] = evalvarn(v);
filling the coefs set the coefs gel(P,i+2)
renormalize P = RgX_renormalize_lg(P, d+3);
The last step will take care of setting the sign correctly.
```

```
t_POL example
```

```
GEN fun(long d, long v)
{
    long i;
    GEN P = cgetg(d+3, t_POL);
    P[1] = evalvarn(v);
    for (i = 0; i <= n; i++)
        gel(P, 2+i) = sqrs(i);
    return RgX_renormalize_lg(P, d+3);
}</pre>
```

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t STR

t_STR are character string. GSTR(x) return the string pointer. GEN strtoGENstr(const char \star s) convert a C string to t_STR The number of long to allocate for n characters is nchar2nlong.

t_CLOSURE

 $\ensuremath{\texttt{t_CLOSURE}}$ holds GP functions. The length can be 6,7 or 8.

- 6 inline closure
- 7 function
- 8 true closure

closure_arity(C): arity of the closure.

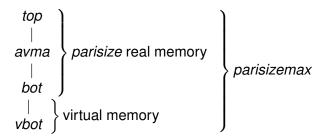
True closures are GP functions that have a non empty context of execution:

Inline closure is code that appear inside loop:

print (i^2+1) is an inline closure (that depend on *i*).

The PARI stack

Since GEN can be quite complex, PARI uses a dedicated memory management system: the PARI stack. The PARI stack is a contiguous chunk of memory used as a scratchpad for computation. It is made of two consecutive chunks (allocated with mmap). The first chunk is of length parisize starts from top down to bot and is allocated as real memory. The second chunk starts from bot down to vbot and is allocated as virtual memory. The total length from top to vbot is parisizemax. The stack pointer is called avma.



When avma reaches bot, the bot is lowered (and a Warning: increasing stack size occurs), When bot reaches vbot, a PARI stack overflow error occurs. The virtual memory between the old and new bot is then converted to real memory.

The low-level function for allocating memory is very simple:

```
INLINE GEN
new_chunk(size_t x) /* x is a number of longs */
{
   GEN z;
   if (x > (avma-bot) / sizeof(long))
        new_chunk_resize(x);
   z = ((GEN) avma) - x;
   avma = (pari_sp)z;
   return z;
}
```

The PARI stack has several advantage.

- memory allocation are very fast.
- it is fully reentrant.
- it prevents memory leak.
- it is always obvious who owns a particular address.
- it allows object to be serialized.

In principle, GEN can exist anywhere in memory, however all libpari functions that return new GENs allocate them on the PARI stack.

A function should normally start by recording the stack pointer avma of type pari_sp and restore the stack at the end. For that purpose, gc_GEN, gc_long, gc_ulong are available.

```
<TYPE> fun(...)
{
    pari_sp av = avma;
    <TYPE> z;
    ...
    z = ...;
    return gc_<TYPE>(av, z);
}
```

where <TYPE> can be any of long, ulong, GEN. If the GEN is known to be a leaf type, gc_leaf should be used. For void function, use set_avma(av).

gc_GEN and gc_upto

gc_GEN (av, z) works by copying recursively the GEN z outsize the stack, reseting avma to av and recopying z at avma. The cost only depend on the size of z gc_upto(av, z) is a faster version that just move z to avma, shiftint the pointers as needed. However it has two requirements.

- 1. the pointer $\ensuremath{\mathrm{z}}$ must be created before its components.
- 2. The part of the stack used by z and its components need to be connected.

GEN produced by gc_GEN always have this property. If furthermore, there were no temporaries created, return z is sufficient.

Examples

```
pari_sp av = avma;
GEN a = utoi(3), b = utoi(4);
GEN V = cgetg(3,t_VEC);
gel(V,1) = a;
gen(V,2) = b;
return gc_GEN(av, V);
```

In this example, the first condition is not respected, gel(V, 1) and gel(V, 2) are created before *V*.

```
GEN V = cgetg(3,t_VEC);
gel(V,1) = utoi(3);
gen(V,2) = utoi(4);
return V;
```

In this example, there is no temporaries created, no need for ${\tt gc}.$

```
pari_sp av = avma;
GEN V = cgetg(3,t_VEC);
gel(V,1) = addiu(shifti(gen_1,128),1);
gen(V,2) = utoi(4);
return gc_GEN(av, V);
```

In this example, the second condition is not respected, the object shifti(gen_1, 128) is a temporary is the middle of *V*.

```
pari_sp av = avma;
GEN z = shifti(gen_1,128);
GEN V = cgetg(3,t_VEC);
gel(V,1) = addiu(z,1);
gen(V,2) = utoi(4);
return gc_upto(av, V);
```

In this example, the temporary is created before V, so now both condition hold.

```
pari_sp av = avma;
GEN a = addiu(shifti(gen_1,128), 1);
GEN V = cgetg(3,t_VEC);
gel(V,1) = a;
gen(V,2) = utoi(4);
return gc_GEN(av, V);
```

In this example, gel (V, 1) is created before V.

mkvec2 and retmkvec2

```
pari_sp av = avma;
V = mkvec2(utoi(3), utoi(4));
return gc_GEN(av, V);
```

In this example, the GEN utoi (3) and utoi (4) are created before V.

```
retmkvec2(utoi(3), utoi(4));
```

retmkvec2 is a macro that ensure that cgetg(3,t_VEC) is called before utoi(3) and utoi(4) are evaluated.

```
#define retmkvec2(x,y)\
    do { GEN _v = cgetg(3, t_VEC);\
        gel(_v,1) = (x);\
        gel(_v,2) = (y); return _v; } while(0)
```