# Constructing CM elliptic curves with minimal Galois image 

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From Karim's email, on the 13th of December, 2021:
(1) A talk you could deliver in about 30 minutes related to an implementation project for PARI/GP. Mention whether it is a past implementation, work in progress or planned work. In the latter two cases, it will be helpful to identify bottlenecks and what help could be useful. "No talk" is an allowed answer. :-)

## Computing the Galois image of a non-CM elliptic curve

Fix a number field $F$ with algebraic closure $\bar{F}$, and an elliptic curve $E_{/ F}$ such that End $\overline{\bar{F}}(E) \cong \mathbb{Z}$. Set $E_{\text {tors }}:=E(\bar{F})_{\text {tors }}$ and $E[m]:=E(\bar{F})[m]$ for any $m \in \mathbb{N}$. We have the Galois representation:

$$
\rho_{E}: G_{F}:=\operatorname{Gal}(\bar{F} / F) \rightarrow \operatorname{Aut}_{\mathbb{Z}}\left(E_{\text {tors }}\right) \cong \mathrm{GL}_{2}(\widehat{\mathbb{Z}})
$$

putting together the $\ell$-adic ones $\rho_{E, \ell^{\infty}}: G_{F} \rightarrow \operatorname{Aut}_{\mathbb{Z}}\left(E\left[\ell^{\infty}\right]\right) \cong \mathrm{GL}_{2}\left(\mathbb{Z}_{\ell}\right)$, where $E\left[\ell^{\infty}\right]:=\lim _{n \in \mathbb{N}} E\left[\ell^{n}\right]$.
Serre (1971) The index $\mathscr{I}(E / F):=\mid$ Aut $_{\mathbb{Z}}\left(E_{\text {tors }}\right): \rho_{E}\left(G_{F}\right) \mid$ is finite.
How can we make this theorem effective?
Lombardo (2015) $\mathscr{I}(E / F)<\exp \left(1.9 \cdot 10^{10}\right) \cdot\left(d_{F} \cdot \max \left\{1, h(E), \log \left(d_{F}\right)\right\}\right)^{12395}$, where $d_{F}:=[F: \mathbb{Q}]$ and $h(E)$ denotes the stable Faltings height of $E$, computed in PARI/GP by ellheight ( $E$ ).
Zywina (2015) The set $\left\{\mathscr{I}(E / \mathbb{Q}): E_{/ \mathbb{Q}\}} \subseteq \mathbb{N}\right.$ is finite, if we assume Serre's uniformity conjecture.
Brau-Avila (2015) There is a very slow, deterministic algorithm which computes an integer $m \geq 2$ such that $\rho_{E}\left(G_{F}\right)=\pi_{m}^{-1}\left(\rho_{E, m}\left(G_{F}\right)\right)$, where $\pi_{m}: \operatorname{Aut}_{\mathbb{Z}}\left(E_{\text {tors }}\right) \rightarrow \operatorname{Aut}_{\mathbb{Z}}(E[m])$ is the reduction map, and $\rho_{E, m}:=\pi_{m} \circ \rho_{E}$. Moreover, there is another algorithm which computes $\rho_{E, m}\left(G_{F}\right)$, and thus can be used to compute $\mathscr{I}(E / F)$.

Rouse, Sutherland \& Zureick-Brown (2021): If $F=\mathbb{Q}$, there is an algorithm, more efficient than Brau-Avila's, which computes $\rho_{E, \ell^{\infty}}\left(G_{F}\right)$, for any prime $\ell \in \mathbb{N}$, using work of Lozano-Robledo (2018).
This algorithm works also without CM, is implemented in MAGMA, and was applied to the 238764310 elliptic curves appearing in the database by Balakrishnan, Ho, Kaplan, Spicer, Stein \& Weigandt (2016).
and this allows us to compute $\mathscr{I}(E / F)$. Indeed, such an $L$ always exists, and one can take $L=F(E[I])$ for any ideal $I \subseteq \mathscr{O}$ such that $|\mathbb{Z} /(I \cap \mathbb{Z})|>\max \left(2,\left|\mathscr{O}^{\times}\right| / 2\right)$. This gives an algorithm to compute $\mathscr{I}(E / F)$.
Challenge: Implement this algorithm completely in PARI/GP.
Fix a number field $F$, and an elliptic curve $E_{/ F}$ such that $\operatorname{End}_{F}(E) \cong \mathscr{O} \subseteq K \subseteq F$, where $\mathscr{O}$ is an order inside an imaginary quadratic field $K$. Then $\rho_{E}\left(G_{F}\right) \subseteq \operatorname{Aut}_{\mathscr{O}}\left(E_{\text {tors }}\right) \cong \widehat{\mathscr{O}}^{\times}$, and $\mathscr{I}(E / F):=\mid$ Aut $_{\mathscr{O}}\left(E_{\text {tors }}\right): \widehat{\mathscr{O}}^{\times} \mid$is finite.
Campagna \& P. (2021) $\mathscr{I}(E / F)=\left[F \cap K^{\mathrm{ab}}: H_{\mathscr{O}}\right] \cdot\left(\left|\mathscr{O}^{\times}\right| /\left[F\left(E_{\text {tors }}\right): F K^{\mathrm{ab}}\right]\right)$, where $H_{\mathscr{O}}=K(j(E))$ is the ring class field of $\mathscr{O}$. Moreover, for any finite $L \supseteq F$ such that $F\left(E_{\text {tors }}\right)=L K^{\text {ab }}$, one has:

$$
\mathscr{I}(E / F)=\frac{\left|\mathscr{O}^{\times}\right| \cdot\left[L \cap K^{\mathrm{ab}}: K\right]}{|\operatorname{Pic}(\mathscr{O})| \cdot[L: F]}
$$ (2016).

## How to face our first challenge

To implement the algorithm, we need the following steps:

- write a function ellcmdisc(E) which takes as input an elliptic curve $E$ defined over a number field $F$ and returns 1 if $E$ does not have complex multiplication, and otherwise returns the discriminant $\Delta_{\mathscr{O}} \in \mathbb{Z}$ of the imaginary quadratic order $\mathscr{O}$ such that $\operatorname{End}_{\bar{F}}(E) \cong \mathscr{O}$;
- write a function elldivfield ( $E, n$ ) which takes as input an elliptic curve $E$ defined over a number field $F$, and an integer $n \in \mathbb{N}$ (or, more generally, an element $n \in \operatorname{End}_{F}(E)$ ), and outputs an irreducible polynomial $f(x) \in K[x]$ such that the number field $F[x] /(f)$ is the $n$-division field $F(E[n])$ of $E$;
- write a function nfrelativize (f,g) which, given two irreducible polynomials $f, g \in \mathbb{Q}[x]$ and some embedding $K:=\mathbb{Q}[x] /(f) \hookrightarrow F:=\mathbb{Q}[x] /(g)$, returns a polynomial $h(y) \in K[y]$ such that $F \cong K[y] /(h)$;
- write a function subab (K,f,\{flag $=0\}$ ) which, given a number field $K$ and an irreducible polynomial $f \in K[x]$, returns another irreducible polynomial $g \in K[x]$ such that $K[x] /(g)$ is the maximal abelian sub-extension of $L:=K[x] /(f)$. If $f$ lag $=1$, return just the degree $\left[L \cap K^{\mathrm{ab}}: K\right]$.
Then, we can easily compute $\mathscr{I}(E / F)$, using quadclassunit to compute $|\operatorname{Pic}(\mathscr{O})|$.

Fix a number field $F$ and an elliptic curve $E_{/ F}$ such that $\operatorname{End}_{F}(E) \cong \mathscr{O} \subseteq K \subseteq F$.
Another good choice for a finite extension $L \supseteq F$ such that $F\left(E_{\text {tors }}\right)=L K^{\mathrm{ab}}$, is given by $L=F(\sqrt[u]{\alpha})$, where $u=\left|\mathscr{O}^{\times}\right|$and $\alpha \in F$ is such that $E$ is the twist by $\sqrt[u]{\alpha}$ of an elliptic curve $E_{/ F}^{\prime}$ such that $F\left(E_{\text {tors }}^{\prime}\right)=F K^{\text {ab }}$.
Campagna \& P. (2020) For any $\mathscr{O}$ such that $\Delta_{\mathscr{O}} \notin\left\{-4 f^{2} \mid f=p_{1}^{a_{1}} \cdots p_{r}^{a_{r}}, p_{1} \equiv \cdots \equiv p_{r} \equiv 1(4)\right\}$, there are infinitely many non-isomorphic elliptic curves $E_{/ H_{\mathscr{O}}}^{\prime}$ such that $H_{\mathscr{O}}\left(E_{\text {tors }}^{\prime}\right)=K^{\mathrm{ab}}$.
The previous theorem gives an algorithm to construct such an $E^{\prime}$. More precisely (if $\Delta_{\mathscr{O}}<-4$ and $K \neq \mathbb{Q}(i)$ ):

- take any elliptic curve $E_{/ H_{\mathscr{O}}}$ such that End $_{H_{\mathscr{O}}}(E) \cong \mathscr{O}$, e.g. $E=\operatorname{ellfromj}(j(\mathscr{O}))$;
- if $H_{\mathscr{O}}(E[3]) \subseteq K^{\mathrm{ab}}$, then $\mathscr{I}\left(E / H_{\mathscr{O}}\right)=2$, and we can take $E^{\prime}=E$;
- if $H_{\mathscr{O}}(E[3]) \nsubseteq K^{\mathrm{ab}}$, continue as follows:
- take any (e.g. the smallest) prime $p \in \mathbb{N}$ which splits in $K$, is inert in $\mathbb{Q}(i)$, and such that $p \nmid f_{\mathscr{C}} \cdot N_{H_{\mathscr{O}} / \mathbb{Q}}\left(f_{E}\right)$;
- find $\alpha \in H_{\mathscr{O}}$ such that $H_{\mathscr{O}}(E[\mathfrak{p}])=H_{\mathfrak{p}, \mathscr{O}}(\sqrt{\alpha})$, where $p \cdot \mathscr{O}=\mathfrak{p p}$ and $H_{p, \mathscr{O}}$ is the $p$-th ray class field of $\mathscr{O}$;
- take $E^{\prime}=E^{(\alpha)}$.

Challenge: Implement this algorithm completely in PARI/GP.
Theoretical challenge: Find $A, B \in \mathbb{Q}(u, v)$ such that $E^{\prime}: y^{2}=x^{3}+A\left(j(\mathscr{O}), \Delta_{\mathscr{O}}\right) \cdot x+B\left(j(\mathscr{O}), \Delta_{\mathscr{O}}\right)$ for each $\mathscr{O}$.

$$
\alpha:=13956546560 \cdot(1190435+2307955 i-1032149 \sqrt{-5}+532379 \sqrt{5})
$$

and $E^{\prime}=E^{(\alpha)}$, which gives $E^{\prime}: y^{2}-\left(\frac{1-i+\sqrt{-5}+\sqrt{5}}{2}\right) x y-\left(\frac{1+i+\sqrt{-5}+\sqrt{5}}{2}\right) y=x^{3}+x^{2}+(2 i-\sqrt{5}) x-1+2 i$.
Consider $\mathscr{O}=\mathbb{Z}[\sqrt{-5}]$, so that $\Delta_{\mathscr{O}}=-20$ and $H_{\mathscr{O}}=\mathbb{Q}(\sqrt{-5}, i)$. Then, we can take:

$$
E: y^{2}=x^{3}+29736960(36023 \sqrt{5}-80550) x-55826186240(16154216 \sqrt{5}+36121925)
$$

so that $j(E)=282880 \sqrt{5}+632000$ and $\mathscr{I}\left(E / H_{\mathscr{O}}\right)=1$. We can take $p=3$. Then, $3 \cdot \mathscr{O}=\mathfrak{p} \bar{p}$ with $\mathfrak{p}=(3, \sqrt{-5}+1)$. Hence, we have the equality $H_{\mathfrak{p}, \mathscr{C}}=H_{\mathscr{O}}$, which is reflected by the factorization:

$$
\begin{aligned}
\phi_{E, 3}(x)=3 \cdot & (x+594880+59840 i-26048 \sqrt{-5}+266816 \sqrt{5}) . \\
& (x+594880-59840 i+26048 \sqrt{-5}+266816 \sqrt{5}) . \\
& \left(x^{2}-(1189760+533632 \sqrt{5}) x-2668089262080-1193205432320 \sqrt{5}\right)
\end{aligned}
$$

of the 3 -division polynomial of $E$. Thus, we have $H_{\mathscr{O}}(E[\mathfrak{p}])=H_{\mathscr{O}}(\sqrt{\alpha})$, where:

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A minimal and a maximal index...
Michelangelo Buonarroti,
La creazione di Adamo

